# Optimal transport networks in spatial equilibrium

Fajgelbaum & Schaal (Econometrica 2020)

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### **Motivation:**

- in theory: transport infrastructure, through trade costs, shapes spatial distribution of prices, real incomes, and trade flows (Allen-Arkolakis 2014; Donaldson-Hornbeck 2016)
- in practice: billions spent on roads, airports, seaports, ...

#### Question: How should we allocate transport infrastructure across space?

- are observed networks suboptimal?
- if so, how large are the losses?

What we need: a method to identify the best set of infrastructure investments

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Application: European road network

- welfare gain from optimal 50% expansion  $\approx$  2.5%
- added links broadly similar to TEN-T proposal



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- 1. local consumption feasibility
- 2. local factor market clearing
- 3. road construction feasibility
- 4. balanced flows
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- optimal allocation sub-problem: given roads and flows, solve neoclassical trade model
  - arbitrary geography, productivity, pop.
  - roads & flows pin down trade costs
  - has well-known solutions

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- optimal flow sub-problem: given roads, what are optimal trade flows?
  - shut off GE effects ⇒ well-known problem from optimal transport lit. (see Galichon 2016, Ch8)
  - well-known solution methods: duality techniques to derive *potential field*
  - here: same methods work in GE in field of prices

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- network design problem:
  - under assumptions: whole problem is convex, can use dual techniques in price space
  - otherwise: use dual techniques to solve sub-problems jointly, then *simulated annealing* to approximate global max

# Model

### **Environment:**

- discrete set of locations  $\mathcal{J} = \{1, \dots, J\}$
- N traded goods, aggregated by  $D_j(\cdot)$
- 1 non-traded good in fixed supply,  $H_j$
- labor  $L_j$  (fixed or mobile)

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Preferences: homothetic and concave

$$u_j = U(c_j, h_j)$$

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- fixed factors  $\mathbf{V}_j = (V_j^1, \dots, V_j^M)$ 
  - immobile across locations
  - mobile across sectors
- intermediates  $\mathbf{X}_j = (X_j^1, \dots, X_j^N)$
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### Nests neoclassical trade models:

- Armington
- Ricardian
- factor-endowment

# Locations ${\cal J}$ are arranged on an undirected graph

- J nodes in set  ${\mathcal J}$
- $\mathcal{E} \equiv$  set of edges (unordered pairs of  $\mathcal{J}$ )
- $\mathcal{N}(j) \equiv$  set of neighbors of location j
- goods only shipped through connected locations  $k \in \mathcal{N}(j)$
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- can handle topography, too: mountains, rivers, ...



### Transport technology: Endogenous iceberg trade costs

Transporting one unit of good n from j to k requires

$$\tau_{jk}^n = \tau_{jk}(Q_{jk}^n, I_{jk})$$

units of the good itself, where

- $Q_{ik}^n \equiv$  flow of good *n* along *jk* link regardless of where it was produced
- $I_{jk} \equiv$  level of infrastructure on jk link

We assume that, along all links jk,



and  $\tau_{jk}(Q, I) \neq \tau_{i\ell}(Q, I)$  because of topography (distance, ruggedness, ...)

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# Infrastructure technology

- Define transport network as  $\{I_{jk}\}_{j \in \mathcal{J}, k \in \mathcal{N}_j}$
- Infrastructure built with one scarce resource: "asphalt"
  - in fixed aggregate supply K
  - freely shipped across locations
  - only used for building infrastructure
  - all together: opportunity cost of building  $I_{jk}$  is just not building  $I_{i\ell}$

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- $\delta_{ik}^{l} \equiv$  asphalt intensity of building infrastructure on link jk
- **network-building constraint**: with multiplier  $P_K$  ("price of asphalt")

$$\sum_{j \in \mathcal{J}} \sum_{k \in \mathcal{N}(j)} \delta_{jk}^{\prime} I_{jk} \leq K$$

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$$\sum_{j\in\mathcal{J}}\sum_{k\in\mathcal{N}(j)}\delta'_{jk}I_{jk}\leq K$$

• problem allows for bounds on infrastructure:

$$0 \leq \underbrace{\underline{I}_{jk}}_{\text{existing road}} \leq I_{jk} \leq \underbrace{\overline{I}_{jk}}_{\text{feasible road}} \leq \infty$$

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- asphalt K is scarce but sunk



## Planner's problem (with fixed labor)

$$W = \max_{\substack{c_j, h_j, \{I_{jk}\}, \\ \{D_j^n, L_j^n, \mathbf{V}_j^n, \mathbf{X}_j^n, \{Q_{jk}^n\}_k\}_n}} \sum_j \omega_j L_j U(c_j, h_j)$$

subject to

- 1. availability of commodities:  $c_j L_j \leq D_j (D_j^1, \ldots, D_j^n)$  and  $h_j L_j \leq H_j$  for all j
- 2. conservation of flows constraint: for all j, n

$$D_j^n + \sum_{n'} X_j^{nn'} + \sum_{k \in \mathcal{N}(j)} (1 + au_{jk}^n) Q_{jk}^n \leq Y_j^n + \sum_{i \in \mathcal{N}(j)} Q_{ij}^n$$

3. **network-building** constraint (subject to bounds on  $I_{jk}$ ):

$$\sum_{j\in\mathcal{J}}\sum_{k\in\mathcal{N}(j)}\delta'_{jk}I_{jk}\leq K$$

- 4. local factor market clearing  $(L_j \text{ and } \mathbf{V}_j)$
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### Convexity of the planner's problem

(i) Given the network  $\{I_{jk}\}$ , the joint optimal transport and allocation problem (with fixed labor) is a **convex optimization problem** if

```
Q\tau_{jk}(Q, I) is convex in Q, \forall j \in \mathcal{J}, k \in \mathcal{N}(j);
```

(ii) the full planner's problem including the network-design problem is a convex optimization problem if, in addition,

 $Q\tau_{jk}(Q, I)$  is convex in Q and I,  $\forall j \in \mathcal{J}, k \in \mathcal{N}(j)$ ;

Intuition: two complementary forces...

- large flows  $\implies$  congestion
- large flows  $\implies$  incentive to increase capacity

#### Eventually, congestion must win out for the full problem to be convex

If the full planner's problem is convex, then the KKT conditions are **sufficient**, and thus the solution to the **dual problem** 

 $\inf_{\mathbf{P}\geq 0}\sup_{\mathbf{x}}\mathcal{L}(\mathbf{x},\mathbf{P}),$ 

coincides with the solution to the primal problem

 $\sup_{\mathbf{x}} \inf_{\mathbf{P} \geq 0} \mathcal{L}(\mathbf{x}, \mathbf{P}).$ 

(i) use FOCs to express control variables as function of prices,  $\mathbf{x}(\mathbf{P})$ 

(ii) reduce dual to convex minimization problem just in space of prices:

 $\inf_{\mathbf{P} \geq 0} \mathcal{L}(\mathbf{x}(\mathbf{P}), \mathbf{P})$ 

(iii) solve with gradient-descent algorithm ightarrow guaranteed to converge to global optimum

Get a no-arbitrage condition from the planner's first-order condition for flows:

$$rac{P_k^n}{P_j^n} \leq 1 + au_{jk}^n + rac{\partial au_{jk}^n}{\partial Q_{jk}^n} Q_{jk}^n, \quad = \ ext{if} \ Q_{jk}^n > 0.$$

Any price differential that exceeds the marginal transport cost would be exploited.

- with no congestion  $(\partial \tau^n_{ik}/\partial Q^n_{ik}=0)$ , price differential bounded by trade cost
- if  $Q\tau_{jk}(Q, I)$  is convex in Q, can invert **RHS** to get  $\tilde{Q}_{jk}^n(P_k^n/P_j^n)$ , which is increasing
- each good *n* flows in only **one direction** (to the higher price)





# Workhorse parameterization: Log-linear transport costs

Suppose that

$$au_{jk}(\boldsymbol{Q},\boldsymbol{I}) = \delta_{jk}^{ au} rac{\boldsymbol{Q}^{eta}}{\boldsymbol{I}^{\gamma}}, \quad ext{with } eta \geq \mathbf{0}, \gamma \geq \mathbf{0}$$

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Also get intuitive, closed-form solutions:

• optimal flows:

$$Q_{jk}^n = \left[rac{1}{1+eta}rac{I_{jk}^\gamma}{\delta_{jk}^ au}\max\left\{rac{P_k^n}{P_j^n}-1,0
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• optimal network:

$$I_{jk} = \min\{\max\{I_{jk}^*, \underline{I}_{jk}\}, \overline{I}_{jk}\}, \text{ with } I_{jk}^* = \left[\frac{\gamma}{P_{\kappa}} \frac{\delta_{jk}^{\tau}}{\delta_{jk}^I} \left(\sum_n P_j^n (Q_{jk}^n)^{1+\beta}\right)\right]^{\frac{1}{1+\gamma}}$$

### Parameterization

- preferences:  $c_j^{\alpha} h_j^{1-\alpha}$ , CES
- production:  $Y_j^n = z_j^n L_j^n$
- transport: log-linear (previous slide)
- traded sectors: 11 (10 diff., 1 homog.)

Vars/params	Calibration
$\mathcal{J}$	0.5 arc-degree cells
ε	connect contiguous cells
$I_{ik}^{obs}$	weighted sum of lanes
$\tilde{L}_{i}^{obs}$	NASA-SEDAC's GPW v.4
$GDP_i^{obs}$	Yale's G-Econ 4.0
$z_j^n$	model-implied
$\tilde{H}_{j}$	model-implied
$\alpha = 0.4$	exp. share of local cons.
eta= 0.13	Couture et al. (2018)
$\gamma=$ 0.10	Couture et al. (2018)
$\delta^{\tau}_{ik}$	distance <sub>jk</sub>
$\delta^{I}_{jk}$	distance <sub>jk</sub> , ruggedness <sub>jk</sub>

### Optimal road expansion and reallocation in Spain



Same exercise across 24 countries: avg. welfare gain of 1.8% for both

# Optimal road expansion across Western Europe vs. TEN-T proposal



Model (labor mobile within countries)

#### total welfare gains of 2.5%



Discretized TEN-T network

Appendix

Full planner's problem can be decentralized with proper set of Pigouvian taxes:

- atomistic shippers pay tolls
- link-specific contractors build roads, earn toll revenue
- tolls set to perfectly offset congestion

### But that set of Pigouvian taxes is not in the DGP!

• the DGP may be a competitive equilibrium, but it's not this one

So why do it this way? Anything weaker and we lose global optimality